

# Scenario Generation and Sampling Methods

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SVAN 2016 – IMPA  
May 11th, 2016

Surveys in Operations Research and Management Science 19 (2014) 56–85



ELSEVIER

Contents lists available at ScienceDirect

## Surveys in Operations Research and Management Science

journal homepage: [www.elsevier.com/locate/sorms](http://www.elsevier.com/locate/sorms)



### Review

## Monte Carlo sampling-based methods for stochastic optimization



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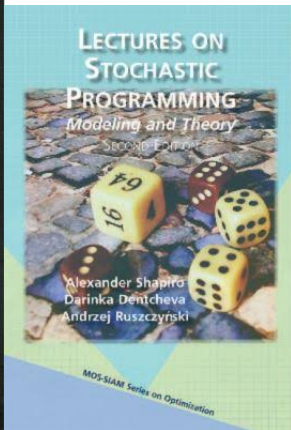
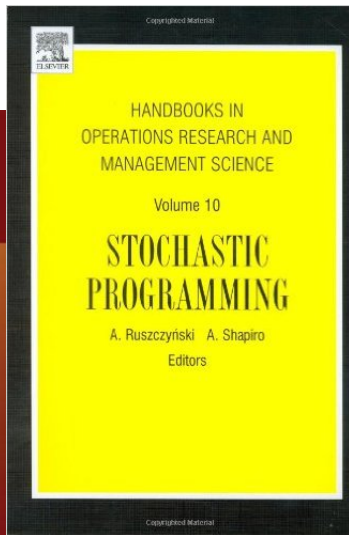
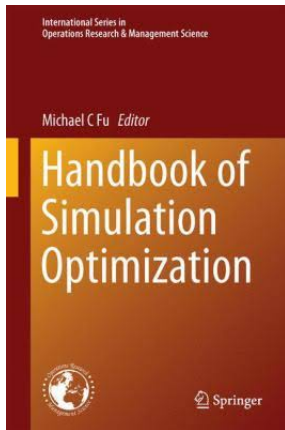
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### HIGHLIGHTS

- We survey the use of Monte Carlo sampling-based methods for stochastic optimization.
- We provide over 240 references from both the optimization and simulation literature.
- We discuss the convergence of optimal solutions/values for sampling approximations.
- Topics related to the implementation of sampling-based algorithms are discussed.
- An overview of alternative sampling techniques to reduce variance is presented.

# Support material



# The SAA approach

Recap: we were studying what happens when we approximate the problem

$$\min_{x \in X} \{g(x) := \mathbb{E}[G(x, \xi)]\} \quad (\text{SP})$$

by

$$\min_{x \in X} \left\{ \hat{g}_N(x) := \frac{1}{N} \sum_{j=1}^N G(x, \xi^j) \right\}. \quad (\text{SP}_N)$$

This is called the Sample Average Approximation (SAA) approach.

# Asymptotic properties of SAA

Let

$\hat{x}_N :=$  an optimal solution of  $(SP_N)$

$S_N :=$  the set of optimal solutions of  $(SP_N)$

$\nu_N :=$  the optimal value of  $(SP_N)$

and

$x^* :=$  an optimal solution of  $(SP)$

$S^* :=$  the set of optimal solutions of  $(SP)$

$\nu^* :=$  the optimal value of  $(SP)$

As the sample size  $N$  goes to infinity, does

- $\hat{x}_N$  converge to some  $x^*$ ?
- $S_N$  converge to the set  $S^*$ ?
- $\nu_N$  converge to  $\nu^*$ ?

# Problems with stochastic constraints

So far our analysis has focused on the problem

$$\min_{x \in X} \{g(x) := \mathbb{E}[G(x, \xi)]\} \quad (\text{SP})$$

which has a *deterministic* feasibility set  $X$ , say,  $X = \{x : h_i(x) \leq 0\}$ ,  $i = 1, \dots, m$ .

**Issue:** What if we have *stochastic* constraints? How to model the problem?

We will consider two classes of problems:

- ①  $X$  has the form  $\mathbb{E}[H_i(x, \xi)] \leq 0$ ,  $i = 1, \dots, m$ .
- ②  $X$  has the form  $P(H_i(x, \xi) \leq 0) \geq 1 - \alpha$ ,  $i = 1, \dots, m$   
(or, equivalently,  $p(x) := P(H_i(x, \xi) > 0) \leq \alpha$ ).

# Problems with expectation constraints

Let us consider the case where the stochastic constraint is  $\mathbb{E}[H(x, \xi)] \leq 0$ .

A natural approach is use SAA and replace the constraint with

$$\frac{1}{N} \sum_{j=1}^N H(x, \xi^j) \leq 0.$$

Does that work?

- Let us consider a simple example where the objective function is  $g(x) = x$ , and the constraint function is  $H(x, \xi) = \xi - x$  where  $\xi$  has distribution  $\text{Normal}(0, \sigma)$ , i.e., the constraint is  $x \geq 0 = \mathbb{E}[\xi]$ .
- The SAA of this constraint is

$$x \geq \frac{1}{N} \sum_{j=1}^N \xi^j.$$

so the SAA solution is  $\hat{x}_N = \frac{1}{N} \sum_{j=1}^N \xi^j$ .

# Problems with expectation constraints

Note that  $\frac{1}{N} \sum_{j=1}^N \xi^j$  has distribution  $\text{Normal}(0, \sigma/\sqrt{N})$ .

So, there is a *50% chance* that the solution  $\hat{x}_N$  will be **infeasible** for the original problem!

**Idea:** Perturb the feasibility set, writing it as

$$U^\varepsilon := \{x \in X : \mathbb{E}[H(x, \xi)] \leq \varepsilon\}.$$

- When  $\varepsilon > 0$  we have a *relaxation* of the original problem.
- When  $\varepsilon < 0$  we have a *tightening* of the original problem.



# Problems with expectation constraints

Now let  $U_N^0$  denote the feasibility region if the SAA problem, i.e.,

$$U_N^0 = \left\{ x \in X : \frac{1}{N} \sum_{j=1}^N H(x, \xi^j) \leq 0 \right\}.$$

## Theorem

*When  $X$  is compact, the function  $H(\cdot, \xi)$  is Lipschitz and  $H(x, \cdot)$  has finite moment generating function, there exist constants  $M$  and  $\beta > 0$  such that*

$$P(U^{-\varepsilon} \subseteq U_N^0 \subseteq U^{\varepsilon}) \geq 1 - Me^{-\beta\varepsilon^2 N}.$$

# Application: Problems with CVaR constraints

Given a random variable  $Z$ , the conditional value-at-risk (CVaR) of  $Z$  is defined as

$$\text{CVaR}_{1-\alpha}[Z] = \frac{1}{\alpha} \int_{1-\alpha}^1 \text{VaR}_{\gamma}[Z] d\gamma$$

where

$$\text{VaR}_{\gamma}[Z] := \min\{t \mid P(Z \leq t) \geq \gamma\}.$$

It is well known that the CVaR can be written as

$$\text{CVaR}_{1-\alpha}[Z] = \min_{\eta \in \mathbb{R}} \left\{ \eta + \frac{1}{\alpha} \mathbb{E}[(Z - \eta)_+] \right\},$$

where  $(a)_+ := \max(a, 0)$ .

Also, the optimal solution  $\eta^*$  of this problem is  $\text{VaR}_{1-\alpha}[Z]$ !

# Connection with CVaR constraints

Note also that, when  $Z$  has **continuous** distribution, we have

$$\begin{aligned}
 \mathbb{E}[Z \mid Z > \text{VaR}_{1-\alpha}[Z]] &= \mathbb{E}[\text{VaR}_{1-\alpha}[Z] + (Z - \text{VaR}_{1-\alpha}[Z]) \mid Z > \text{VaR}_{1-\alpha}[Z]] \\
 &= \text{VaR}_{1-\alpha}[Z] + \frac{\mathbb{E}[(Z - \text{VaR}_{1-\alpha}[Z])_+]}{P(Z > \text{VaR}_{1-\alpha}[Z])} \\
 &= \eta^* + \frac{1}{\alpha} \mathbb{E}[(Z - \eta^*)_+] \\
 &= \min_{\eta \in \mathbb{R}} \left\{ \eta + \frac{1}{\alpha} \mathbb{E}[(Z - \eta)_+] \right\} \\
 &= \text{CVaR}_{1-\alpha}[Z].
 \end{aligned}$$

In particular, this implies that  $\text{CVaR}_{1-\alpha}[Z] \geq \text{VaR}_{1-\alpha}[Z]$ .

# Application: Problems with CVaR constraints

Consider now the problem

$$\begin{aligned} \min_{x \in X} \quad & g(x) \\ \text{s.t.} \quad & \text{CVaR}_{1-\alpha}[F(x, \xi)] \leq a. \end{aligned}$$

Then, by using the optimization representation of CVaR we can write the problem as

$$\begin{aligned} \min_{x \in X, \eta \in \mathbb{R}} \quad & g(x) \\ \text{s.t.} \quad & \eta + \frac{1}{\alpha} \mathbb{E} [(F(x, \xi) - \eta)_+] \leq a, \end{aligned}$$

which falls into the standard formulation by defining

$$H((x, \eta), \xi) := \eta + \frac{1}{\alpha} (F(x, \xi) - \eta)_+ - a.$$

# Chance-constrained problems

Chance constraints can be very helpful in modeling some situations.

This is true especially when what matters is *whether or not* a constraint was violated, not the *amount* of violation. For example,

- Reliability problems
- Problems with physical constraints

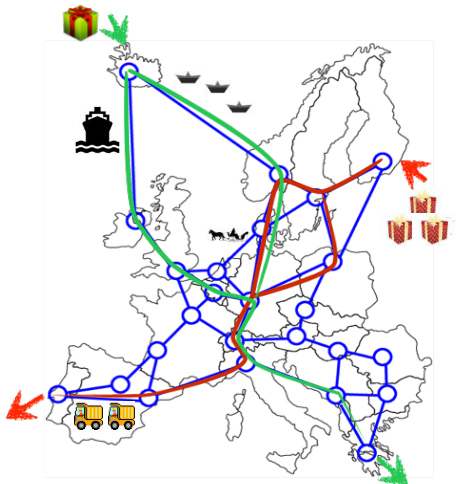
Also, it is often easier to choose the chance constraint level than to choose, say, penalties for violation.

# Example: a telecommunication problem

Network:  $G=(V,A)$

Commodities:  $\mathcal{C}$ , each one with **possible** demand  $d_c$  to be routed from  $s_c$  to  $t_c$

Capacities:  $w_l$  for each link  
(need to be integral)

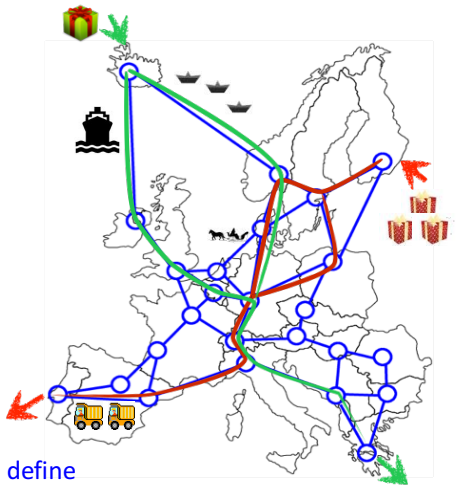


# Example: a telecommunication problem

Network:  $G=(V,A)$

Commodities:  $\mathcal{C}$ , each one with **possible** demand  $d_c$  to be routed from  $s_c$  to  $t_c$

Capacities:  $w_l$  for each link  
(need to be integral)



**Problem:** To **route** each commodity and **define capacities** for each link that minimize the capacity installation cost, subject to a **reliability** constraint

# Modeling the problem

- Each connection  $c$  communicates with probability  $\rho_c$   
( $\xi_c \sim \text{Bernoulli}(\rho_c)$ )
- We need to determine the minimum capacity  $w_\ell$  for each link  $\ell$  that will meet communication requirements with probability at least  $1 - \alpha_\ell$ .
- The routing variables  $x_\ell^c$  are equal to one if connection  $c$  uses link  $\ell$ , zero otherwise



# A chance-constrained formulation

$$\min_{x,w} \sum_{\ell \in L} w_{\ell}$$

s.t.

$$\sum_{\ell \in \delta^+(s_c)} x_{\ell}^c - \sum_{\ell \in \delta^-(s_c)} x_{\ell}^c = -1 \quad \forall c \in \mathcal{C}$$

$$\sum_{\ell \in \delta^+(t_c)} x_{\ell}^c - \sum_{\ell \in \delta^-(t_c)} x_{\ell}^c = 1 \quad \forall c \in \mathcal{C}$$

$$\sum_{\ell \in \delta^+(n)} x_{\ell}^c - \sum_{\ell \in \delta^-(n)} x_{\ell}^c = 0 \quad \forall c \in \mathcal{C}, \forall n \neq s_c, t_c$$

$$P \left( \sum_{c \in \mathcal{C}} \xi_c x_{\ell}^c \leq w_{\ell} \right) \geq 1 - \alpha_{\ell} \quad \forall \ell \in L.$$

# Feasible regions: an example

$$\begin{array}{ll} \min_{x \in \mathbb{R}^2} & c_1 x_1 + c_2 x_2 \\ \text{s.t.} & P(\xi x_1 + x_2 \geq 7) \geq 1 - \alpha \end{array}$$

Assuming  $\xi \sim U[0, 1]$ , draw the feasible region  $C(\alpha)$  for  $\alpha = 0.3$  and  $\alpha = 0.7$ .

# Solution

If  $\xi \sim U[0, 1]$  then the feasible set is

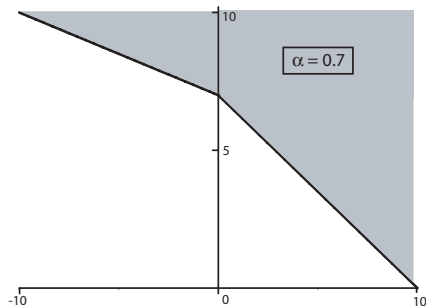
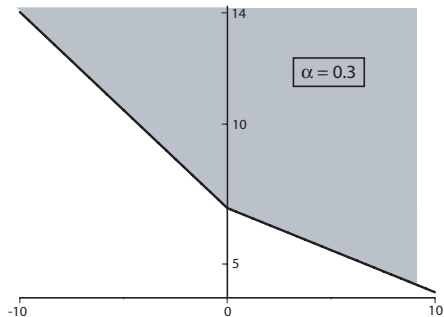
$C(\alpha) = C_+(\alpha) \cup C_0(\alpha) \cup C_-(\alpha), \alpha \in (0, 1)$ , where

$$C_+(\alpha) = \{x \in \mathbb{R}^2 \mid x_1 > 0, \alpha x_1 + x_2 \geq 7\},$$

$$C_0(\alpha) = \{(0, x_2) \in \mathbb{R}^2 \mid x_2 \geq 7\},$$

$$C_-(\alpha) = \{x \in \mathbb{R}^2 \mid x_1 < 0, (1 - \alpha)x_1 + x_2 \geq 7\}$$

# Feasible regions



# Connection with CVaR constraints

Consider the chance constraint

$$P(F(x, \xi) \leq 0) \geq 1 - \alpha.$$

Note that this is equivalent to

$$\text{VaR}_{1-\alpha}[F(x, \xi)] \leq 0.$$

Recall that we saw earlier that  $\text{CVaR}_{1-\alpha}[Z] \geq \text{VaR}_{1-\alpha}[Z]$ .

- Therefore, if we replace the chance constraint  $P(F(x, \xi) \leq 0) \geq 1 - \alpha$  with  $\text{CVaR}_{1-\alpha}[F(x, \xi)] \leq 0$ , we have a *conservative approximation*.
- The advantage of such an approximation is that the feasibility set is convex if  $F(\cdot, \xi)$  is convex.

# Sampling approaches

- Non-convexity of chance-constraints does not occur when the distribution of  $\xi$  belongs to a certain class (called log-concave distributions).
- But what to do if the random parameters do not follow a tractable distribution?
- One alternative is to apply the SAA approach, which replaces the chance constraint by its sample average.
- The resulting problem is easier to solve, and provides useful information to the true problem.

## SAA

- Let  $\xi^1, \dots, \xi^N$  be a random sample from  $\xi$ .
- Using that  $P(\xi \in A) = \mathbb{E}[\mathbb{I}_A(\xi)]$ , the SAA of a chance constrained problem is

$$\begin{aligned} & \min_{x \in X} g(x) \\ \text{s.t. } & p_N(x) := \frac{1}{N} \sum_{j=1}^N \mathbb{I}_{(0, \infty)}(H(x, \xi^j)) \leq \gamma \end{aligned}$$

(Compare with the original problem:)

$$\begin{aligned} & \min_{x \in X} g(x) \\ \text{s.t. } & p(x) := P(H(x, \xi) > 0) \leq \alpha \end{aligned}$$

# The scenario approach

Note that if we take  $\gamma = 0$  in the above formulation we obtain

$$\begin{aligned} \min_{x \in X} \quad & g(x) \\ \text{s.t.} \quad & H(x, \xi^j) \leq 0, \quad j = 1, \dots, N. \end{aligned}$$

If each function  $H(\cdot, \xi)$  is *convex* and  $g$  is *convex*, then the above problem is convex.

- This is called the **scenario approach**.
- What is the relationship to the original problem?



# The scenario approach

## Theorem

Select a confidence parameter  $\beta \in (0, 1)$ , and let  $d_x$  denote the dimension of  $x$ . Suppose that  $H(\cdot, \xi)$  is convex. If

$$N \geq \frac{2}{\alpha} \left( \ln \frac{1}{\beta} + d_x \right),$$

then, with probability at least  $1 - \beta$  we have that  $\hat{x}_N$  satisfy all constraints in the original problem but at most a fraction  $\alpha$ , that is,

$$P(H(\hat{x}_N, \xi) > 0) \leq \alpha,$$

regardless of the distribution of  $\xi$ .

# An equivalent IP formulation

Consider now the case of  $\gamma > 0$ .

- Given a sample of size  $N$ , we can write the problem as

$$\begin{aligned} & \min_{x \in X} g(x) \\ \text{s.t. } & H(x, \hat{\xi}^i) - \mathcal{M}z_i \leq 0 \quad i = 1, \dots, N, \\ & \frac{1}{N} \sum_{i=1}^N z_i \leq \gamma, \\ & z_i \in \{0, 1\}^N. \end{aligned} \tag{P}$$

- That is, we obtain an IP formulation, which is particularly helpful when  $H$  is linear in  $x$ .

# Feasibility results

Similar results to the scenario approach theorem (i.e., feasibility of  $\hat{x}_N$  guaranteed up to a confidence  $1 - \beta$ ) can be obtained, under various different settings:

- When  $X$  is finite;
- When  $H(x, \xi)$  is of the form  $H(x, \xi) = \xi - h(x)$ ;
- When  $H(\cdot, \xi)$  is a Lipschitz function, with Lipschitz constant independent of  $\xi$ .

# Asymptotic results

Condition (A): There is an optimal solution  $\bar{x}$  of the true problem such that for any  $\epsilon > 0$  there is  $x \in X$  with  $\|x - \bar{x}\| \leq \epsilon$  and  $p(x) < \alpha$ .

## Consistency of SAA

Suppose that

- (i) the significance levels of the true and SAA problems are the same, i.e.,  $\gamma = \alpha$ ,
- (ii) the set  $X$  is compact,
- (iii) the function  $g(x)$  is continuous,
- (iv)  $H(x, \xi)$  is a Carathéodory function,
- (v) condition (A) holds.

Then,  $\nu_N \rightarrow \nu^*$  and  $\text{dist}(\hat{S}_N, S^*) \rightarrow 0$  w.p.1 as  $N \rightarrow \infty$ .

# Dealing with small probabilities

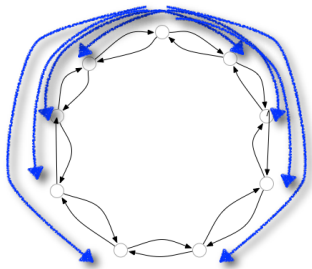
Let us consider again the reliability problem seen earlier, and suppose the reliability factor is very small, say,  $10^{-6}$ .

What happens to the SAA approximation?

- As we saw earlier, the sample size estimates to achieve some desirable confidence are proportional to  $1/\alpha$ .
  - This is not surprising: the probability that the first violation occurs in the  $k$ th sample is  $(1 - \alpha)^{k-1}\alpha$ .
  - Therefore, on average we need  $(1 - \alpha)/\alpha$  samples just to obtain one case for which violation occurs!
- So, we need **a lot** of samples.
- But each sample corresponds to a variable in the IP formulation!

# Why not just use $\alpha = 0$ ?

$\alpha = 0$  : Shortest Path Solution  
(optimal)



$\rho = 0.1$

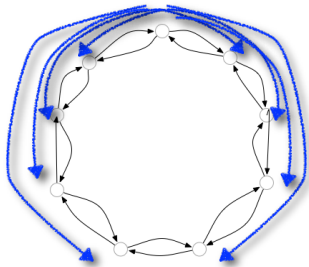
No. of connections routed on each link = 10

Capacity  $w_l$  on each link = 10

Total cost:  $18 \times 10 = 180$

# Why not just use $\alpha = 0$ ?

$\alpha = 0$  : Shortest Path Solution  
(optimal)

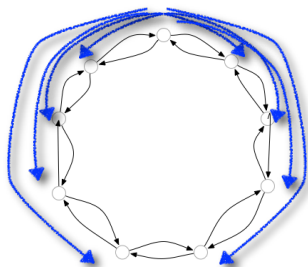


$\rho = 0.1$

No. of connections routed on each link = 10  
Capacity  $w_l$  on each link = 10

Total cost:  $18 \times 10 = 180$

$\alpha = 10^{-6}$  : Shortest Path Solution



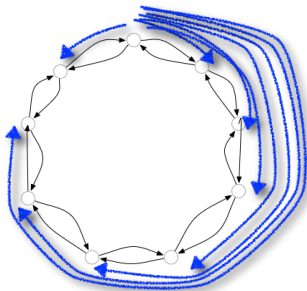
$\rho = 0.1$

No. of connections routed on each link = 10  
Capacity  $w_l$  on each link = 7

Total cost:  $18 \times 7 = 126$

# Why not just use $\alpha = 0$ ?

$\alpha = 10^{-6}$ : Optimal Solution



$\rho = 0.1$

No. of connections routed on each clockwise link = 28

No. of connections routed on each counterclockwise link = 1

Capacity  $w_l$  on each clockwise link = 12

Capacity  $w_l$  on each c/clockwise link = 1

Total cost:  $9 \times 12 + 9 \times 1 = 117$



# Lessons from this example

- We cannot pretend that a very small  $\alpha$  is equivalent to zero...
- On the other hand, when  $\alpha$  is very small SAA will require a lot of samples!
- We need to do some "smarter sampling"
- One such strategy is **importance sampling** — see Guzin's talk!